

# Abstractions for Dynamics of Automata Networks (brief overview)

Loïc Paulevé

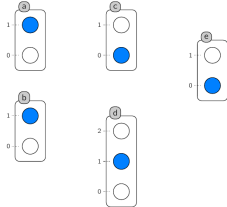
CNRS/LRI, Université Paris-Sud, Orsay, France – BioInfo team

`loic.pauleve@lri.fr`

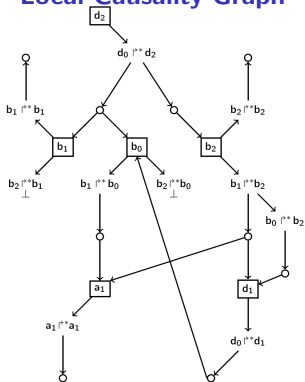
`http://loicpauleve.name`

November 23, 2015 - Journée nationale du groupe de travail BLOSS, Paris

Automata network

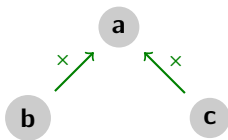


Local Causality Graph



Over-approximation of reachability  
 Under-approximation of reachability  
 Under-approximation of cut sets  
**Goal-oriented model reduction**

## Transition-centered specification

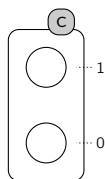
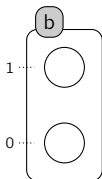
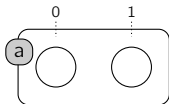


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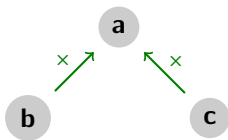
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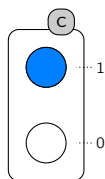
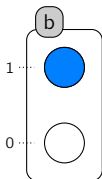
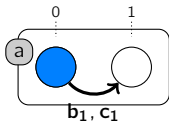


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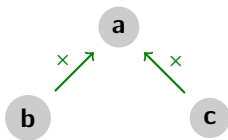
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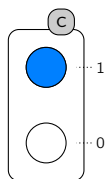
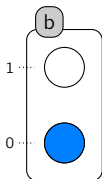
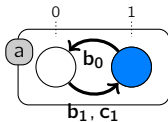


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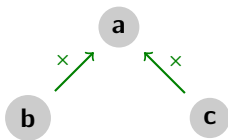
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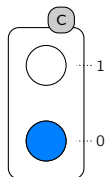
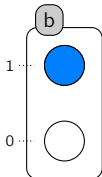
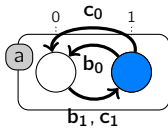


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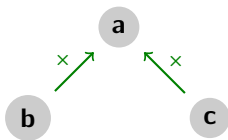
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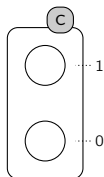
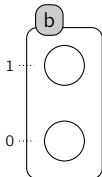
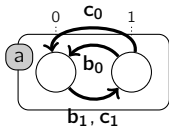
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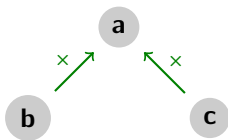
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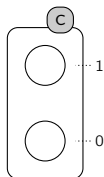
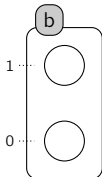
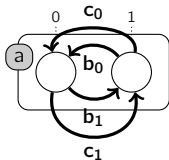
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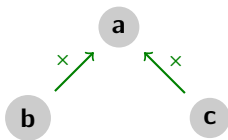
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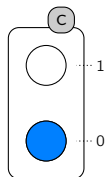
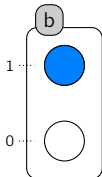
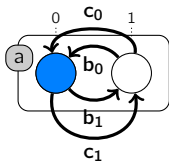
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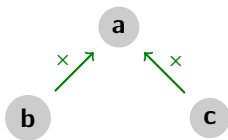
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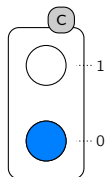
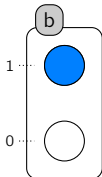
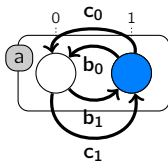
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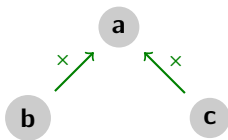
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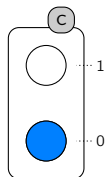
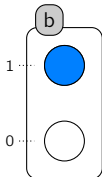
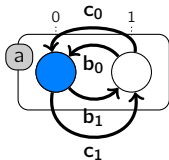
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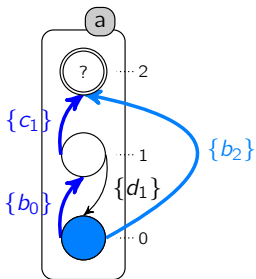
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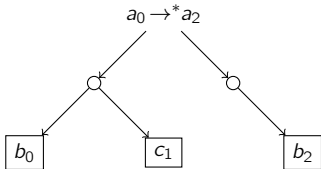


## Local Causality

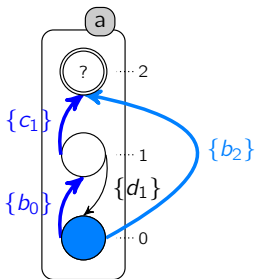


$$\text{local-cause}(a_0 \rightarrow^* a_2) = \{a_0 \xrightarrow{b_0} a_1 \xrightarrow{c_1} a_2, \\ a_0 \xrightarrow{b_2} a_2\}$$

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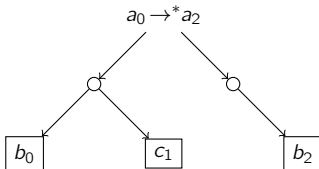


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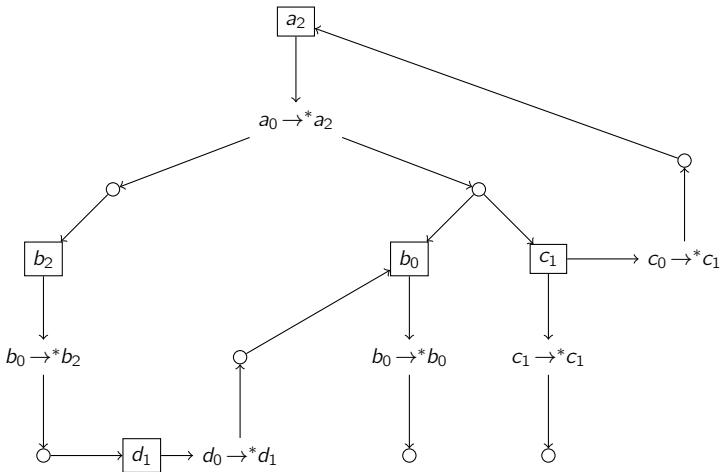


For any trace  $\pi$  starting at some global state  $s$  with  $a_0 \in s$  and reaching  $a_2$ :

- either  $a_0 \xrightarrow{b_0} a_1 \xrightarrow{c_1} a_2$  or  $a_0 \xrightarrow{b_2} a_2$  is a sub-trace of  $\pi$ ;
- either  $b_1$  and  $c_0$ , or  $b_2$  are reached before  $a_2$  in  $\pi$ .

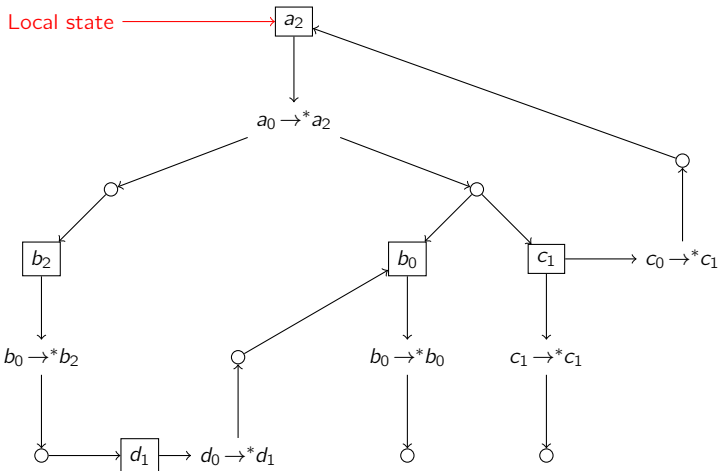
## Local Causality Graph

- Causality of  $a_2$ .
- Initial context  $\varsigma = \{a \mapsto \{0\}; b \mapsto \{0\}; c \mapsto \{0, 1\}; d \mapsto \{0\}\}$ .



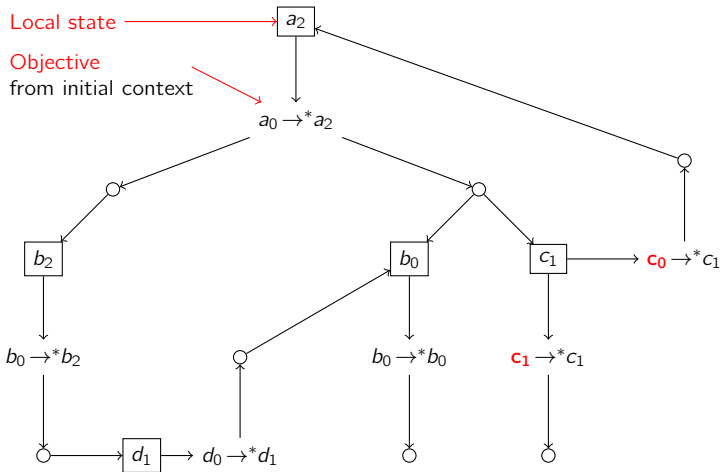
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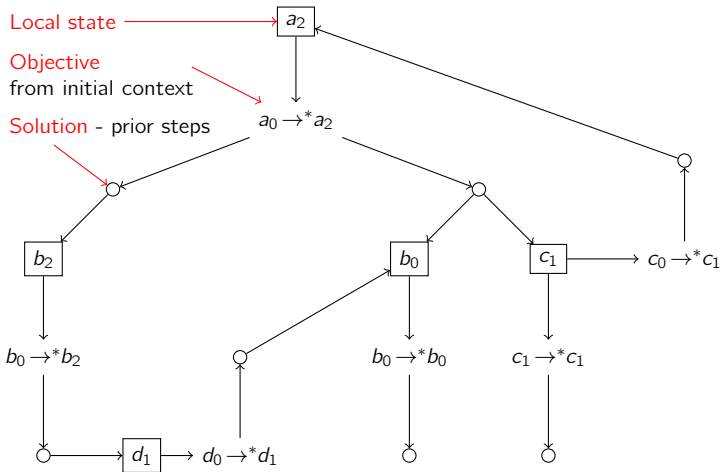
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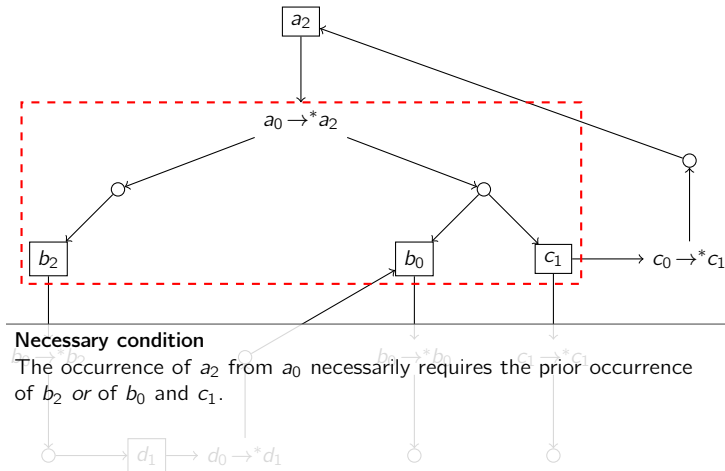
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## Results using local causality analysis

### Verification for transient reachability

- Necessary or sufficient conditions;
- Low complexity (polynomial with nb automata, exp with automaton size);

### Cut-sets for transition reachability (towards control)

- Identify sets of mutations to break a transition reachability;
- Under-approximation of cut-sets
- Applied to models with 10,000 nodes.

### Goal-oriented model reduction

- Given a reachability property (goal), identify **“useless” transitions**.
- Preserves **all minimal traces** for the goal reachability.
- Low complexity (polynomial with nb automata, exp with automaton size);

⇒ pre-processing before model-checking

⇒ benchmarks show effective model reduction

(research report: <https://hal.archives-ouvertes.fr/hal-01149118>)

## Software: Pint

<http://loicpauleve.name/pint>



- Input: automata networks
  - implemented in LogicalModels (shipped with GINsim)
- Command line tools:
  - Static analysis for reachability
  - “Cut sets” for reachability
  - Model reduction w.r.t. reachability property
  - Fixed points
  - Inference of Thomas parameters
  - Interface with model-checkers (NuSMV, ITS, mole).
- OCaml library (C/C++ bindings in progress)

(secret project: GUI)

## Conclusion/Perspectives

### Static analysis of automata network dynamics

- Local **causality analysis**, summarized in a compact graph
- Reasoning on **necessary/sufficient conditions** for transitions
- Multiple applications for **transient reachability** analysis
- Can address very large networks (several 1,000 nodes).

### Goal-oriented model reduction

- Embed in model-checking approaches (Petri net unfoldings, etc.)
- Embed in **model identification** (static pruning of model space).

### Cell reprogramming

- Large models
- Boolean abstraction well suited (experimental validations)
- Causality analysis for **deriving potential reprogramming determinants?**

Thanks for your attention!