Coarse-Grained Brownian Dynamics Simulation of Rule-Based Models

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Abstract. Studying spatial effects in signal transduction, such as colocalization along scaffold molecules, comes at a cost of complexity. In this paper, we propose a coarse-grained, particle-based spatial simulator, suited for large signal transduction models. Our approach is to combine the particle-based reaction and diffusion method, and (non-spatial) rulebased modeling: the location of each molecular complex is abstracted by a spheric particle, while its internal structure in terms of a site-graph is maintained explicit; The particles diffuse inside the cellular compartment and the colliding complexes stochastically interact according to a rulebased scheme. Since rules operate over molecular motifs (instead of full complexes), the rule set compactly describes a combinatorial or even infinite number of reactions. The method is tested on a model of Mitogen Activated Protein Kinase (MAPK) cascade of yeast pheromone response signaling. Results demonstrate that the molecules of the MAPK cascade co-localize along scaffold molecules, while the scaffold binds to a plasma membrane bound upstream component, localizing the whole signaling complex to the plasma membrane. Especially we show, how rings stabilize the resulting molecular complexes and derive the effective dissociation rate constant for it.

1 Introduction

Signal transduction pathways can contain several proteins and activation steps which give rise to complex spatiotemporal dynamics. Interactions between signaling molecules do not only transmit activations, but can also localize the molecules to certain structures or compartments in the cell [20]. This work aims at a simulator that is able to handle the complexity, and includes the localization of the molecules. Current modeling techniques include ordinary differential equations (ODE), or, if space is important, the corresponding partial differential equations (PDE). If the stochasticity is significant, the chemical or reaction-diffusion master equation (CME/RDME) describes the system dynamics. More specifically, each molecule can be tracked individually in the simulation, e.g. using Brownian/Smoluchowski or Green's function dynamics (see e.g. review [13]).

Transient complex formation of proteins and their post-translational modification in signaling can lead to a combinatorial number of distinct molecular species. Rule-based languages, such as Kappa [7] or BioNetGen [9], provide a compact representation of such combinatorial processes. Nonspatial rule-based models can be efficiently simulated [3], and are amenable to further quantitative analysis, e.g. formal model reduction [10, 11].

An extended Kappa framework was proposed in order to model the internal spatial structure of complexes that form during the process [4]. The Meredys simulator [21] in addition includes the molecule positions. Furthermore SRSim [12] provides a high resolution spatial rule-based extension, tracking the position and internal structure of all complexes. Indeed, the possibility to specify the binding angles sometimes naturally enforces a unique assembly path of a desired complex structure: for example, a polymer chain with angles between the bonds of $\pi - 2\pi/N$ in a plane will form a ring of N monomers, and local rules are sufficient to describe the global structure [4]. However, the exact molecular geometry of signaling molecules is often not known. In such cases, the simulation including the binding angles becomes more complex, without contributing additional insights.

Here we present a framework that supports a particle-based, spatial simulation, but omits the internal geometry. Still, as the internal structure in terms of a site-graph is maintained explicit, their cooperative effect on complex stability can be investigated effectively. In particular, we derive the dissociation rate constant for rings. The simulation is applied to MAPK (Mitogen activated protein kinase) signaling in yeast, where both localization and activation is mediated by a scaffold [20].

The paper is structured as follows. In Sect. 2, the general particle-based framework and the biophysical principles of complex formation are introduced. The formal framework underlying the simulator is outlined in Sect. 3. In Sect. 4, we show the application to MAPK signaling and we discuss the results.

2 Coarse-grained Particle Diffusion and Reaction

The mobility of molecules in the cytoplasm is mainly governed by diffusion. Diffusion can be modeled efficiently by a random walk of the molecules of interest such that the myriad of solvent molecules can be omitted. We assume that the properties of the solvent allow us to use the Stokes-Einstein equation to obtain the (translational) diffusion coefficient $D_i \propto r_i^{-1}$ for a given molecular radius r_i of particle *i*. We also assume that the actual shape of the molecules averages to a sphere with radius r_i at the temporal resolution of the method. Therefore only the position but not the orientation of the molecules has to be tracked.

2.1 Remark on Diffusion-Controlled Reactions

Molecules can only react with each other if they are in contact/collide, and the collision process is governed by diffusion. Accordingly the observable bulk/macroscopic reaction rate constant k_{ij} between particles *i* and *j* in solution is determined both by the rate constant of collisions $k_D(i, j) = 4\pi(r_i + r_j)(D_i + D_j)$ (in 3D space) and the microscopic rate constant k'_{ij} that determines the reactive fraction of collisions [2]. For spherically symmetric molecules with isotropic reaction properties the microscopic and macroscopic rate constant are related in the form $k_{ij}^{-1} = k_D(i,j)^{-1} + k'_{ij}^{-1}$ (Collins-Kimball model). Reactions with very high $k_{ij} \approx k_D(i,j)$ require $k'_{ij} \to \infty$, i.e. every collision leads to a reaction such that they are diffusion-controlled. In contrast, reactions with $k_{ij} \ll k_D(i,j)$ are reaction-controlled; in this case $k'_{ij} \approx k_{ij}$. A diffusion factor/function can be introduced $f_D(i,j) := k_D(i,j)/(k_D(i,j) + k'_{ij})$ such that $k = f_D k'$ [2].

More in detail, signaling molecules have specific reaction sites, i.e. nonisotropic reaction properties. Such molecules have to be in contact (by translational diffusion) and correctly aligned with their reaction sites (by rotational diffusion). We define the corresponding nanoscopic rate constant as k'' and the conversion factor f_{DR} such that $k = f_{DR}k''$, however the derivation of f_{DR} is not straightforward and several approximations exist [2, 22]. For completeness we also introduce $f_R = f_{DR}f_D^{-1}$ for the conversion $k' = f_Rk''$, i.e. for integrating out the rotational diffusion effect only. In general: $k \leq k' \leq k''$.

Reversible reactions $A + B \rightleftharpoons C$ with forward reaction at rate constant k_{AB} and backward k_C require to scale both rate constants with f(A, B) in order to maintain the macroscopic reaction equilibrium (dissociation constant $K_d = k_C/k_{AB} = k'_C/k'_{AB} = k''_C/k''_{AB}$) [17, 14]. For consistency we define the respective conversion factors f = 1 for unimolecular reactions that are not reversed by a bimolecular reaction, such that microscopic rate constants are always defined.

2.2 General Particle-Based Diffusion and Reaction Method

The present method implements the λ - ρ model [8], i.e. a discrete time continuous space random walk for diffusion. Each particle position x_i is updated by

$$\boldsymbol{x}_{\boldsymbol{i}}(t + \Delta t) = \boldsymbol{x}_{\boldsymbol{i}}(t) + \sqrt{2D_{\boldsymbol{i}}\Delta t}\boldsymbol{\xi}$$
(1)

with diffusion coefficient D_i and standard normal random variable $\boldsymbol{\xi}$. Particles can overlap as discussed in [14], so collision testing is only needed for (static) reaction compartment boundaries. Reactions are executed with a probability that depends on their arity as follows.

Unimolecular reactions $A \to \ldots$ at rate constant k_A are executed in this method in every step with probability

$$P_A = 1 - exp(-k'_A \Delta t) \approx k'_A \Delta t \text{ if } \Delta t \to 0$$
⁽²⁾

for each molecule which is of type A [8]. Note that the Bernoulli-trial scheme leads to a binomial distribution, which converges to the Poisson distribution for small probabilities in each step (law of rare events). The Poissonian reaction process has exponentially distributed inter event waiting times as expected.

Bimolecular reactions $A + B \rightarrow \ldots$ can only occur if two molecules are closer than their collision distance $(r_A + r_B)$. If so, then the reaction is executed with probability

$$P_{A,B} = \frac{k'_{AB}\Delta t}{4\pi (r_A + r_B)^3/3}$$
(3)

as derived in [16]. The accuracy constraint $P_{A,B} < 0.2$ gives an upper bound for Δt , if larger Δt are needed the reaction probability of [8] has to be used.

Higher order and Hill-type reaction schemes: for a particle-based simulation higher order reaction models have to be composed into their elementary uni- and bimolecular reaction steps (in which they also occur in nature).

2.3 Complexes

Complexes C that form out of $A + B \rightarrow C$ can still be modeled as spherical particles. The complex radius r_C is obtained e.g. under the assumption that the volume/mass of A and B is redistributed into C with constant mass density, i.e.

$$r_C = \left(\frac{3}{4\pi} \left[\frac{4\pi r_A^3}{3} + \frac{4\pi r_B^3}{3}\right]\right)^{1/3} = \left(\sum_{i \in C} r_i^3\right)^{1/3} \tag{4}$$

as suggested in [23]. Alternatives for Eq. (4) are discussed and listed in Table 1 in [13]. The diffusion coefficient D_C is given by $D_C = D_0/r_0r_C^{-1}$ based on a reference D_0 and r_0 in the Stokes-Einstein relation. Reactions between complexes α and β can occur when they are within their contact distance. The joint $k'_{\alpha\beta} = \sum_{A \in \alpha} \sum_{B \in \beta} k'_{AB}$ (assume $k'_{AB} = 0$ if no reaction between A and B is defined) could directly be used to calculate the binding probability by Eq. (3). However we decided to track each reaction individually as described in Sect. 3. Therefore also all resulting bonds in a complex are tracked individually. Such a bond in complex α between $A \in \alpha$ and $B \in \alpha$ can break with the corresponding dissociation rate constant/probability defined for the interaction. But if the two formerly directly will stay together – and **aligned with their binding sites**. Therefore A and B can rebind in every step with P_{AB} , leading to bond recovery with rate constant

$$\bar{k}_{AB} = \frac{3k''_{AB}}{4\pi(r_A + r_B)^3} \tag{5}$$

effectively in a first order reaction (cf. Eq. 3 and Eq. 2) [15]. The recovery reaction does not take place in the (relatively large) reaction volume of the whole reaction compartment but in the (relatively tiny) interaction volume of the two agents. Therefore \bar{k} is large compared to all other rates, and the lifetime of the open state of the bond $\tau \propto \bar{k}^{-1}$ is accordingly relatively small (cf. Test Case 2).

3 Spatial Stochastic Simulation of Rule-based Model

Each particle of the simulation is an instantiation of a molecular species. A molecular species can be a protein, its post-translationally modified form or a protein complex that consists of proteins bound together. In order to reflect this internal structure of molecular species we represent them by *site-graphs*, in which modifications of protein residues and bonds are explicitly encoded, as introduced in Kappa [7].

Notations. We denote by $\{e_1 \mapsto v_1, \dots, e_N \mapsto v_N\}$ the mapping from distinct elements e_i to values v_i . Given a mapping A, dom(A) denotes its domain, and A(e) the value associated to e in A. We also write $A\{e \mapsto v\}$ for the mapping A updated so that e maps to v; and $A \setminus e$ for the mapping A updated so that the mapping from e is removed.

3.1 Site-graphs

A site-graph is an undirected graph where typed nodes have sites, and edges are partial matchings on sites. Moreover, the sites which do not serve for forming edges are called *internal*, and they are assigned a value from a predefined set. The nodes of the site-graph can be interpreted as protein names, and sites of a node stand for protein binding domains. Internal states are used to encode post-translational modifications.

Let S denote the set of site labels, and \mathcal{I} the set of internal values that can be assigned to sites. The function $I : S \to \mathcal{P}(\mathcal{I})$ denotes the set of internal values that a site s can take. Let \mathcal{A} be the set of node types. Each node type is being equipped with a set of sites, defined by a signature map $\Sigma : \mathcal{A} \to \mathcal{P}(S)$. Finally, the set of admissible bindings between sites is defined by the mapping $\mathcal{E} : \mathcal{A} \times S \to \mathcal{P}(\mathcal{A} \times S)$ so that if $(a', s') \in \mathcal{E}(a, s)$ then necessarily $(a, s) \in \mathcal{E}(a', s')$. A rule-based model is defined over a fixed *contact map* defined by the tuple $(\mathcal{A}, \Sigma, \mathcal{E}, I)$ that we consider constant in the rest of this section.

Definition 1. A site-graph is a tuple $G = (V, T, F, E, \psi)$ with

- 1. a set of nodes V,
- 2. a node type function $T: V \to \mathcal{A}$,
- 3. a node interface function $F: V \to \mathcal{P}(\mathcal{S})$, such that for $v \in V$, $F(v) \subseteq \Sigma(T(v))$,
- 4. a set of edges between sites of different nodes, encoded by the function $E: V \times S \to V \times S$ such that if E(v, s) = (v', s') then necessarily $v \neq v'$, E(v', s') = (v, s), and $(T(v'), s') \in \mathcal{E}(T(v), s)$.
- 5. a site evaluation function $\psi: V \times \{s \in \mathcal{S} \mid I(s) \neq \emptyset\} \to \mathcal{I}$, so that $\psi(v, s) \in I(s)$.

Site-graphs will be used in two different contexts: (i) to model physically existing complexes, also termed *concrete site-graphs* (or reaction mixtures), and (ii) to specify the local interaction patterns (rewrite rules). The concrete site-graphs must have all interfaces complete, in the sense that, for all nodes $v \in V$, $F(v) = \Sigma(T(v))$.

Definition 2. Site-graph $G = (V, T, F, E, \psi)$ is a *union* of two site-graphs $G_1 = (V_1, T_1, F_1, E_1, \psi_1)$ and $G_2 = (V_2, T_2, F_2, E_2, \psi_2)$, denoted by $G = G_1 \oplus G_2$, if $V_1 \cap V_2 = \emptyset$, and $V = V_1 \cup V_2$, $F = F_1 \cup F_2$, $E = E_1 \cup E_2$, $\psi = \psi_1 \cup \psi_2$.

Definition 3. Given a site-graph $G = (V, T, F, E, \psi)$, a sequence of edges $(((v_1, s_1), E(v_1, s_1)), \ldots, ((v_k, s_k), E(v_k, s_k)))$ such that for $i = 1, \ldots, k, v_i \in V$,



Fig. 1. Rule application. Rule R_i can be applied to a reaction mixture \mathcal{G} via the embedding σ (indicated by the dotted arrows), and resulting in the reaction mixture \mathcal{G}' . The contact map is such that $\mathcal{A} = \{A, B, C\}, \ \mathcal{S} = \{a, b, c, d, e, f, g\},\$ and $I(a) = \{u, p\}$. More details on the formalism can be found e.g. in [19].

 $s_i \in F(v_i)$, and for i = 1, ..., k - 1, $E(v_i, s_i) = (v, s) \Rightarrow v_{i+1} = v$ and $s_{i+1} \neq s$, is called a *path* between nodes v_1 and v_k . A site-graph G is connected, denoted cc(G) if there exists a path between every two nodes v and v'.

Two site-graphs can be related by an embedding function, which is important for defining the applicability of a rule to a reaction mixture (cf. Fig. 1).

Definition 4. The *embedding* σ between site-graphs $G = (V, T, F, E, \psi)$ and $G' = (V', T', F', E', \psi')$, is induced by a support function $\sigma^* : V \to V'$, if

- 1. σ^* is injective: for all $v, v' \in V$, $[\sigma^*(v) = \sigma^*(v') \implies v = v']$;
- 2. for all $v \in V$, $T(v) = T'(\sigma^*(v))$;
- 3. for all $v \in V$, $[s \in F(v) \implies s \in F'(\sigma^{\star}(v))];$
- 4. $\{(v,s) \mapsto (v',s')\} \subset E \implies \{(\sigma^*(v),s) \mapsto (\sigma^*(v'),s')\} \subset E';$
- 5. $\{(v, s) \mapsto i\} \subset \psi \implies \{(\sigma^*(v), s) \mapsto i\} \subset \psi'$.

If σ^* is bijective, then σ is an *isomorphism*. The set of embeddings between the site-graph G and G' is denoted by embed(G, G'). The set of all embeddings is denoted by \mathbb{E} .

3.2 **Rule-based models**

Definition 5. Consider three types of elementary transformations of site-graphs, denoted by $\delta_{ae}, \delta_{de}, \delta_{ci} : \mathbb{E} \times \mathbb{G} \to \mathbb{G}$, with the following form:

- 1. $\delta_{ae}^{v,s,v',s'}(\sigma,G) = (V,T,F,E\{(\sigma^*(v),s) \mapsto (\sigma^*(v'),s'), (\sigma^*(v'),s') \mapsto (\sigma^*(v),s)\}, \psi)$ (adding an edge) 2. $\delta_{de}^{v,s}(\sigma, G) = (V, T, F, E \setminus E((\sigma^*(v), s)) \setminus (\sigma^*(v), s), \psi)$ (deleting an edge); 3. $\delta_{ci}^{v,s,i}(\sigma, G) = (V, T, F, E, \psi\{(\sigma^*(v), s) \mapsto i\})$ (changing the state value),

where $G = (V, T, F, E, \psi) \in \mathbb{G}, \sigma \in \mathbb{E}$ induced by the suport function $\sigma^* : V_i \mapsto$ V, V_i being a set of nodes such that $v, v' \in V_i$; $v \neq v'$, $(T(v'), s') \in \mathcal{E}(T(v), s)$, and $i \in I(s)$.

A rule R_i is a triple (G, δ, c) , where G is a site-graph, δ is of type δ_{ae} , δ_{de} , or δ_{ci} , and c is a non-negative real number. Applying the rule to a site-graph \mathcal{G} is unique for an embedding $\sigma \in \mathsf{embed}(G, \mathcal{G})$, and results in $\mathcal{G}' = \delta(\sigma, \mathcal{G})$ (for a rigorous explanation, see [5]). In particular, for the identity support function $\sigma^* = \mathbb{I}$, and for $\mathcal{G} = G$, we get $G' = \delta(\mathbb{I}, G)$, which is sometimes called the right-hand-side of the rule (cf. Fig. 1).

3.3 Stochastic Abstract Machine

This subsection defines the syntax and semantics of the abstract formal machine for our coarse-grained spatial stochastic simulation of rule-based models.

In our machine, a *complex* is associated to a connected concrete site-graph, denoted by \mathcal{G} , and a 3-D position. The site-graph of a complex can be modified according to the rules, and may undergo a split into two site-graphs (dissociation), adding a new complex in the machine; or a merge with another sitegraph (association), removing the other complex from the machine. The radius of agents $a \in A$ is determined by the constant r_a , and the radius of the connected site-graph \mathcal{G} is given by $\operatorname{rad}(\mathcal{G})$, according to Eq. (4) (Sect. 2.3):

$$\mathsf{rad}((V, T, F, E, \psi)) = \left(\sum_{v \in V} r_{T(v)}^3\right)^{1/3} \tag{6}$$

The syntax of the abstract machine is given in Def. 6. A machine term M is a quadruple (t, C, X, \mathcal{R}) where t is the current time, C is a map from a complex i to its current site-graph \mathcal{G} , X is a map from a complex i to its position $\mathbf{x} \in \mathbb{R}^3$, and \mathcal{R} the set of rules, as defined in Sect. 3.2.

The execution of the machine at time t follows Eq. (7). First, the position of each complex is updated according to a Brownian diffusion during a fixed Δt time (diffuse function, Def. 7). Second, the active rules at time t are applied according to the probabilities introduced in Sect. 2. The new site-graphs and positions maps give the new term of the machine at time $t + \Delta t$.

$$\begin{aligned}
X' &= \operatorname{diffuse}(X, C, \Delta t) \\
(C', X'') &= \operatorname{react}(C, X', \mathcal{R}, \Delta t) \\
(t, C, X, \mathcal{R}) &\to (t + \Delta t, C', X'', \mathcal{R})
\end{aligned} \tag{7}$$

The react function (Def. 12) applies the rules embeddings that are active at time t in a random order. Each active rule embedding (Def. 11) is specified by the set of concerned complexes (either one or two), and the quadruple (σ, G, δ, k) where (G, δ, k) is a rule in \mathcal{R} , and σ is the actual embedding being a map from the nodes of the rule left hand-side G to the nodes of the concrete complex site-graphs. These rule embeddings gathers the embeddings from any rule in \mathcal{R} to a single complex site-graph C(i) (unary function) or to the union of two complex site-graphs $C(i_1) \oplus C(i_2)$, assuming that complexes i_1 and i_2 have their distance $|X(i_1) - X(i_2)|$ less than the sum of their radii (neighbors function). A rule with embedding σ is actually applied if (1) the embedding is still valid, i.e., previous rules application have not interfered with it; (2) the random number ζ_1 uniformly distributed in [0; 1] is less than the rule probability.

The probability of applying a rule embedding (prob function, Def. 8) depends on the arity of the reaction and on the modification type, as described in Sect. 2. The application of rule embedding to the site-graphs and position maps is given by the do function (Def. 10). In the case of agent site values changes (δ_{ci}) or an internal bound creation (δ_{ae} within one complex *i*), the embedded rule modification is applied to the concerned site-graphs, without any side-effect. In the case

Stochastic Abstract Machine Definitions

Definition 6. Syntax of the abstract machine. i_1, \dots, i_N are the complexes identifiers, assumed all distinct.

$M ::= (t, C, X, \mathcal{R})$	Time t, complex maps C and X , rules \mathcal{R}
$C ::= \{i_1 \mapsto \mathcal{G}_1, \cdots, i_N \mapsto \mathcal{G}_N\}$	Map from a complex i to its site-graph \mathcal{G}
$X ::= \{i_1 \mapsto \mathbf{x}_1, \cdots, i_N \mapsto \mathbf{x}_N\}$	Map from a complex i to its position $\mathbf{x} \in \mathbb{R}^3$

Definition 7. Euler-Maruyama integration of diffusion. Returns the new position maps of complexes after a Brownian diffusion during Δt time, where ζ is a standard normal random variable, and D_0 and r_0 are constants (cf. Eq. (1), Sect. 2.2). Note: state dependent diffusion in different compartments is introduced in Appendix A.

diffuse $(X, C, \Delta t) \stackrel{\Delta}{=} \{i \mapsto | X(i) + \zeta \sqrt{2D_i \Delta t} | D_i = D_0 r_0 / \mathsf{rad}(C(i)), i \in \mathsf{dom}(X)\}$

Definition 8. Rule probability. The rates k', k'' and \bar{k} refers to the rule rate k modified according to Sect. 2, $\operatorname{agents}(\delta_{ae})$ refers to the couple of agents concerned by the bound creation, and r_{a_i} is the radius of agent a_i .

$$\begin{aligned} & \operatorname{prob}(\{i\}, \delta_{\operatorname{ci}}, \sigma, k, C, \Delta t) \stackrel{\Delta}{=} k' \Delta t \\ & \operatorname{prob}(\{i\}, \delta_{\operatorname{de}}, \sigma, k, C, \Delta t) \stackrel{\Delta}{=} k'' \Delta t \text{ if } \operatorname{cc}(\delta_{\operatorname{de}}(\sigma, C(i))) \text{ else } k' \Delta t \\ & \operatorname{prob}(\{i\}, \delta_{\operatorname{ae}}, \sigma, k, C, \Delta t) \stackrel{\Delta}{=} \bar{k} \Delta t \stackrel{\Delta}{=} \frac{3k'' \Delta t}{4\pi (r_{a_1} + r_{a_2})^3} \text{ if } \operatorname{agents}(\delta_{\operatorname{ae}}) = (a_1, a_2) \\ & \operatorname{prob}(\{i_1, i_2\}, \delta, \sigma, k, C, \Delta t) \stackrel{\Delta}{=} \frac{3k' \Delta t}{4\pi (\operatorname{rad}(C(i_1)) + \operatorname{rad}(C(i_2)))^3} \text{ if } i_1 \neq i_2 \end{aligned}$$

Definition 9. Complex formation and dissociation. Returns the modified site-graphs and positions mappings after a merge or a split of complexes. The condition $i_2 \notin dom(C)$ ensures that the new complex i_2 is a fresh identifier.

$$\mathsf{merge}(i_1, i_2, \delta, \sigma, C, X) \stackrel{\Delta}{=} ((C \setminus i_2) \{ i_1 \mapsto \delta(\sigma, C(i_1) \oplus C(i_2)) \}, X \setminus i_2)$$
$$\mathsf{split}(i_1, \delta, \sigma, C, X) \stackrel{\Delta}{=} (C \{ i_1 \mapsto \mathcal{G}_1, i_2 \mapsto \mathcal{G}_2 \}, X \{ i_2 \mapsto X(i_1) \})$$
$$\mathbf{if} \ i_2 \notin \mathsf{dom}(C), \mathcal{G}_1 \oplus \mathcal{G}_2 = \delta(\sigma, C(i_1)), \mathsf{cc}(\mathcal{G}_1), \mathsf{cc}(\mathcal{G}_2)$$

Definition 10. Rule application. Returns the modified site-graphs and positions mappings. Predicate $cc(\mathcal{G})$ is true if and only if \mathcal{G} is connected (Def. 3, Sect. 3.1).

$$\begin{split} &\mathsf{do}(I, \delta_{\mathsf{ci}}, \sigma, C, X) \stackrel{\Delta}{=} (C\{i \mapsto \delta_{\mathsf{ci}}(\sigma, C(i)) \mid i \in I\}, X) \\ &\mathsf{do}(\{i\}, \delta_{\mathsf{ae}}, \sigma, C, X) \stackrel{\Delta}{=} (C\{i \mapsto \delta_{\mathsf{ae}}(\sigma, C(i))\}, X) \\ &\mathsf{do}(\{i_1, i_2\}, \delta_{\mathsf{ae}}, \sigma, C, X) \stackrel{\Delta}{=} \mathsf{merge}(i_1, i_2, \delta_{\mathsf{ae}}, \sigma, C, X) \text{ if } i_1 \neq i_2 \\ &\mathsf{do}(\{i\}, \delta_{\mathsf{de}}, \sigma, C, X) \stackrel{\Delta}{=} (C\{i \mapsto \delta_{\mathsf{de}}(\sigma, C(i))\}, X) \text{ if } \mathsf{cc}(\delta_{\mathsf{de}}(\sigma, C(i)) \\ &\mathsf{do}(\{i\}, \delta_{\mathsf{de}}, \sigma, C, X) \stackrel{\Delta}{=} \mathsf{split}(i, \delta_{\mathsf{de}}, \sigma, C, X) \text{ if } \mathbf{not} \, \mathsf{cc}(\delta_{\mathsf{de}}(\sigma, C(i)) \end{split}$$

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Stochastic Abstract Machine Definitions (Continued)

Definition 11. Active embeddings. unary and binary returns the embedding specifications of rules \mathcal{R} to single or couples of complexes, respectively; neighbors returns the couples of complexes close enough to react; act_embeds returns the (I, E) couples where I is the set of complexes concerned by the embedding specification E.

$$\begin{split} & \mathsf{unary}(i,C,\mathcal{R}) \stackrel{\Delta}{=} \{(\sigma,G,\delta,k) \mid (G,\delta,k) \in \mathcal{R}, \sigma \in \mathsf{embed}(G,C(i))\} \\ & \mathsf{binary}(i_1,i_2,C,\mathcal{R}) \stackrel{\Delta}{=} \{(\sigma,G,\delta,k) \mid (G,\delta,k) \in \mathcal{R}, \sigma \in \mathsf{embed}(G,C(i_1) \oplus C(i_2)), \\ & \sigma \notin \mathsf{embed}(G,C(i_1)) \cup \mathsf{embed}(G,C(i_2))\} \\ & \mathsf{neighbors}(C,X) \stackrel{\Delta}{=} \{\{i_1,i_2\} \mid \mathsf{rad}(C(i_1)) + \mathsf{rad}(C(i_2)) \geq |X(i_1) - X(i_2)|, \\ & i_1,i_2 \in \mathsf{dom}(C), i_1 \neq i_2\} \\ & \mathsf{act_embeds}(C,X,\mathcal{R}) \stackrel{\Delta}{=} \{(\{i\},E) \mid i \in \mathsf{dom}(C), E \in \mathsf{unary}(i,C,\mathcal{R})\} \\ & \cup \{(\{i_1,i_2\},E) \mid \{i_1,i_2\} \in \mathsf{neighbors}(C,X), \\ & E \in \mathsf{binary}(i1,i_2,C,\mathcal{R})\} \end{split}$$

Definition 12. Sequential application of active rule embeddings, order is assumed to be random; returns the modified site-graphs and position maps. The set of concerned complexes is denoted by I, $\bigoplus_{i \in I} C(i)$ denotes the union of site-graphs of complexes in I, ζ_1 is a random variable uniformly distributed within 0 and 1, and embed is the set of embeddings from rule left hand-side to reaction mixture (cf. Sect. 3.1).

apply({
$$(I, (\sigma, G, \delta, k))$$
} $\cup Q, C, X, \Delta t$) $\stackrel{\Delta}{=}$ apply($Q, C, X, \Delta t$)
if $\sigma \notin \text{embed}(G, \oplus_{i \in I} C(i))$
or $\zeta_1 \ge \text{prob}(I, \delta, \sigma, k, C, \Delta t)$
 $\stackrel{\Delta}{=}$ apply($Q, C', X', \Delta t$)
if $\sigma \in \text{embed}(G, \oplus_{i \in I} C(i))$,
 $\zeta_1 < \text{prob}(I, \delta, \sigma, k, C, \Delta t)$,
 $(C', X') = \text{do}(I, \delta, \sigma, C, X)$
apply($\emptyset, C, X, \Delta t$) $\stackrel{\Delta}{=} (C, X)$
react($C, X, \mathcal{R}, \Delta t$) $\stackrel{\Delta}{=}$ apply(act_embeds($C, X, \mathcal{R}), C, X, \Delta t$)



Fig. 2. Comparison of particle-based simulation according to the presented algorithm (black symbols) and ODE model (red line, see Appendix B) for two test cases: (a) enzymatic activation, Test Case 1 and (b) trimerization, Test Case 2. (c) Overview of the reactions in the ring formation process of example (b), binding sites are omitted here.

of a complex formation (δ_{ae} between two complexes i_1 and i_2), the two complexes are merged into one and receives the position of one of the two complexes in a non-deterministic manner (merge function, Def. 9). Finally, in the case of a bond deletion within a complex (δ_{de}), unless the site-graph is still connected, the complex is split in two complexes receiving the same position (split function, Def. 9).

4 Test Cases and Application

Test Cases: The accuracy of the simulation was tested on the principal reaction motifs of signaling models, where we assume the reaction-limited case for simplicity, i.e. k = k' = k''. The test cases are (in Kappa like syntax, see Fig. 1 for agents and sites):

- 1. Reversible enzymatic activation $A(d) + C(b) \rightleftharpoons A(d^1), C(b^1), A(d^1), C(b^1, a_u) \rightarrow A(d^1), C(b^1, a_p)$ and $C(a_p) \rightarrow C(a_u)$.
- 2. Reversible trimerization by: $A(d) + C(b) \rightleftharpoons A(d^1), C(b^1),$ $A(e) + B(g) \rightleftharpoons A(e^2), B(g^2) \text{ and } C(c) + B(f) \rightleftharpoons C(c^3), B(f^3)$

Fig. 2 depicts the simulation result which exactly match the ODE models derived in Appendix B, thus showing the correctness of the approach.

Test Case 2 exemplifies the formation of a ring and the cooperativity of the bonds in the ring. The ABC trimer can only dissociate if two bonds break and are open at the same time. For equal binding rate constants k_1 , equal dissociation rate constants k_2 and under the quasi-steady-state assumption for complexes with one broken bond the effective dissociation rate constant becomes

$$k^* = 6k_2^2 / (\bar{k}_1 + 2k_2) \tag{8}$$

as shown in Appendix B, Eq. (13). In a ring consisting of N agents the effective dissociation rate constant $k^* = (N(N-1)k_2^2) / (\bar{k}_1 + (N-1)k_2)$, will converge to Nk_2 , i.e. with rising N the cooperativity vanishes. Especially small rings are therefore stabilized due to the high bond recovery rate constant \bar{k} . In the present example $k_1 = 5 \times 10^5 M^{-1} s^{-1}$, $k_2 = 0.2 s^{-1}$, $r_A = r_B = r_C = 5 nm$, $D_A = D_B = D_C = 1 \mu m^2/s$. Then $\bar{k}_1 = 198 s^{-1}$, and $k^* = 0.0012 s^{-1} \ll k_2$, i.e. the ring structure is extremely stable. Diffusion limit $k_D = 1.51 \times 10^8 M^{-1} s^{-1} \gg k_1$ i.e. reaction-limited regime.

Signaling with Scaffolds in Space and Time: In order to exemplify how the presented method can be used to model signaling in space and time, we simulate the yeast pheromone response signaling model from Thomson et al. [20] (see Fig. 3a) up to the MAPK Fus3. In the model the actual signaling molecules Ste4, Ste11, Ste7 and Fus3 can bind to the scaffold Ste5. In contrast to Thomson et al. we assume that any activation by the upstream molecule also involves a possible binding interaction between these molecules – although that interaction might be weak. The resulting additional bonds, shown in red in Fig. 3b, enable the formation of three rings around Ste5, which can stabilize



Fig. 3. (a) Yeast mating MAPK signaling pathway from [20]. (b) Binding and activation interactions around the scaffold Ste5. Incomplete signaling complexes cannot transmit the signal and do not include rings, which makes them less stable. (c) Localization of Ste5 to the plasma membrane. (d) Activation of the MAPK cascade. The initial fast activation phase of Fus3 takes place in scaffold based complexes that formed prior to activation. Further activations require that new Ste7-Fus3 pairs form (diffusion-reaction process) which gives rise to slower kinetics. (e) Visualization of the complexes in 3D space at t = 184s using ZigCell3D from ScienceVisuals [24]. The transparent spheres with constant radius of 12.5 nm are used for all complexes. Nucleus and plasma membrane are not shown.

complete signaling complexes. The binding interactions Ste4-Ste11 and Ste11-Ste7 were omitted in [20] because pheromone response must only be activated if Ste5 is present. In the present model the binding interaction was weak enough such that likewise activation requires the presence of Ste5 (data not shown). Effectively, the bond only establishes along the scaffold, such that crosstalk with other signaling pathways can be prevented. The arising rings including Ste5 and the resulting low dissociation due to cooperativity of the bonds in the rings furthermore makes overexpression of Ste5 less harmful than in [20]. This is due to the fact that without the bonds parallel to Ste5, each Ste5-Ligand bond would arise independent of the other ligands. Too many Ste5 instances would therefore lead to complexes where most likely just 1 ligand is connected to Ste5 such that

the signaling cascade is disrupted. In contrast, the additional bonds drive the equilibrium towards complete complexes.

The spatial and temporal dynamics of the activation process in the MAPK signaling cascade is depicted in Fig. 3c and d. Ste4 is a membrane bound protein, therefore all complexes containing Ste4 will likewise be membrane bound (cf. Appendix A). Fig. 3c shows how Ste5 accumulates at the plasma membrane accordingly. The 3D positions of molecules and complexes in the simulation are shown in Fig. 3e.

5 Discussion

The present work shows how a simulation for complex signal transduction models can be implemented, making use of a rule-based model description and including biophysical aspects as well as the spatial component. The present algorithm uses a coarse-grained description and simplified models for binding and dissociation rate constants on the macro-, micro- and nanoscopic level, thus refining the simulation method suggested in Meredys [21]. The applied rate constants and conversion factors as introduced in Sect. 2.1 could be further refined by more detailed models, which e.g. require to solve the reaction-diffusion or corresponding master equation [17, 18].

An accuracy and performance test of the particle-based simulator core in [14] shows that the presented method performs at least as good as Smoldyn [1]. Smoldyn however can handle rule-based models only via the libMoleculizer plugin. Association and dissociation of complexes require costly graph traversal. However, associations require that two complexes are in contact, which is a rare event, and dissociation does not occur more frequently (at steady state). Given the computational cost of the random walk of the molecules and neighbor finding alone, the rule-based extension does not dramatically slow down simulation. The performance of the algorithm can be improved by using a Gillespie scheme (draw exponential waiting time for reaction event) instead of the Bernoulli trials for unimolecular reactions as in [15]. In that case also the order of unimolecular reactions is given by the (ordered) waiting times instead of the random order we proposed in order to execute all reactions in an unbiased manner. For bimolecular reactions it is extremely unlikely that more than two complexes are within the reactive distance such that in most cases there is no ambiguity which reaction is to be applied. Further improvements of the performance could come from multi-scale or mixed approaches for different domains of the simulation [13].

The coarse-grained rate constants enable the calculation of emerging properties like the cooperativity between bonds, that stabilizes rings. In the MAPK signaling example also rings are formed around the scaffold Ste5. Future work can analyze the stability of signaling complexes and the dose-response curve for Ste5 now including the cooperative effect of the bonds. The simulation already includes localization sites that determine the location of the agents (cytoplasmic or membrane based as shown in Appendix A). In the future we are planning to include the transport into the nucleus such that the complete signaling pathway can be analyzed. Furthermore formal model checking of spatial rule-based models has to be included to ensure meaningful models and simulation results.

Additional Information

The appendix is available from http://www.bison.ethz.ch/research/spatial_ simulations_si/CMSB2013_Appendix.pdf and the simulator as well as example files from http://www.bison.ethz.ch/research/spatial_simulations.

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Appendix A: Spatial Rule-Based Model Checking

The following model has to be checked for undefined diffusion states:



A modifiable localization site on agents defines their location and mobility type. The value can be:

- 0 'null': no diffusion (D = 0), i.e. static position ('bound to fixed anchor')
- 1 'pm': 2D diffusion along the plasma membrane ('bound to membrane 1').
- 2 'nm': 2D diffusion along the nuclear membrane ('bound to membrane 2').
- 3 '3D': 3D diffusion in the cytoplasmic domain ('unbound').

However, there can be reachable states with multiple location site values in one complex. In such a case, the smallest location value determines the behavior of the whole complex, and the states 'pm' and 'nm' become conflicting, because a complex can not safely be diffused along both membranes.

In order to prevent such a conflict, one has to ensure that no complex referencing both 'pm' and 'nm' location can be formed. While this condition can be enforced in the abstract machine by assigning a zero the probability to a binding between a complex with 'pm' location and a complex with 'nm' location, it is worth noticing that, in some cases, a conflict-free simulation can be warranted by static analysis ([6]).

Indeed, from neighbors function (Def. 11), two complexes i_1, i_2 can merge only if they are close enough: for any machine term (t, C, X, \mathcal{R}) , $\{i_1, i_2\} \in$ neighbors $(X, C) \Leftrightarrow \operatorname{rad}(C(i_1)) + \operatorname{rad}(C(i_2)) \geq |X(i_1) - X(i_2)|$. Assuming that complex i_1 has location 'pm' and complex i_2 has location 'nm', i_2 and i_2 may be merged only if $\operatorname{rad}(C(i_1)) + \operatorname{rad}(C(i_2)) \geq d$, where d is the minimal distance between nuclear and plasma membranes. It is hence sufficient that the sum of the radius of the largest possible complexes i_1 and i_2 is (strictly) less than d to ensure a conflict-free simulation. The largest complexes can be over-approximated from the contact map of the rule-based model, giving all the possible compositing agents. Note that in some models, it may be possible to have unbounded complex sizes: in such a case, the sufficient condition can not be satisfied, and a dynamical disabling of rules leading to conflict has to be used.

Appendix B: ODE Models for the Test Cases

The following ODEs describe the dynamics of the test cases in the macroscopic, well mixed limit.

Test Case 1: The reversible dimerization of A and B to AB is described by the concentration of the 3 species c_A , c_B and c_{AB} and rates $r_1 = k_1 c_A c_B$ for association and $r_2 = k_2 c_{AB}$ for dissociation. In addition the test case case includes activation from B to B^* if bound to 'enzyme' A with rate constant k_3 and spontaneous deactivation from B^* to B with rate constant k_4 , independent of whether B^* is bound to A. The active form B^* has association and dissociation rate constants k_1^* and k_2^* respectively. This system has 7 rates, the 2 from above and $r_3 = k_1^* c_A c_{B^*}$, $r_4 = k_2^* c_{AB^*}$, $r_5 = k_3 c_{AB}$, $r_6 = k_4 c_{AB^*}$ and $r_7 = k_4 c_{B^*}$, which determine in the ODE system for the 5 species

$$\begin{split} \dot{c}_A &= -r_1 + r_2 - r_3 + r_4 \\ \dot{c}_B &= -r_1 + r_2 & + r_7 \\ \dot{c}_{AB} &= +r_1 - r_2 & - r_5 + r_6 \\ \dot{c}_{AB^*} &= & + r_3 - r_4 + r_5 - r_6 \\ \dot{c}_{B^*} &= & - r_3 + r_4 & - r_7 \end{split}$$

Test Case 2: Symmetric binding of A + B, B + C and C + A all with binding rate constant k_1 and dissociation constant k_2 . The full complex ABC contains three bonds in a ring structure. If a monomer (e.g. A) binds to a dimer (e.g. BC), two bonds will have to be established accordingly, which is described by the joint rate constant $k_3 = 2k_1$. The bond recovery rate constant \bar{k}_1 within the full complex is given by Equation (5). The rates are

$$\begin{array}{ll} r_1 = k_1 c_A c_B & r_2 = k_2 c_{AB} & r_3 = k_1 c_B c_C & r_4 = k_2 c_{BC} \\ r_5 = k_1 c_C c_A & r_6 = k_2 c_{CA} & r_7 = k_3 c_A c_{BC} & r_8 = k_3 c_B c_{CA} \\ r_9 = k_3 c_C c_{AB} & r_{10} = k_2 c_{ABC} & r_{11} = \bar{k}_1 c_{AB.C} & r_{12} = \bar{k}_1 c_{A.BC} \\ r_{13} = \bar{k}_1 c_{ABC} & r_{14} = k_2 c_{AB.C} & r_{15} = k_2 c_{A.BC} & r_{16} = k_2 c_{ABC} , \end{array}$$

where the dot in ABC denotes the location of the one open bond in the ring. Corresponding ODE System (where the multiplier in front of the rates has to be chosen according to the number of bonds that are affected):

$$\dot{c}_{A} = -r_{1} + r_{2} - r_{5} + r_{6} - r_{7} + r_{14} + r_{16}$$

$$\dot{c}_{B} = -r_{1} + r_{2} - r_{3} + r_{4} - r_{8} + r_{14} + r_{15}$$

$$\dot{c}_{C} = +r_{3} - r_{4} - r_{5} + r_{6} - r_{9} + r_{15} + r_{16}$$

$$\dot{c}_{AB} = +r_{1} - r_{2} - r_{9} + r_{15} + r_{16}$$

$$\dot{c}_{BC} = +r_{3} - r_{4} - r_{7} + r_{14} + r_{16}$$

$$\dot{c}_{CA} = +r_{5} - r_{6} - r_{8} + r_{14} + r_{15}$$

$$\dot{c}_{ABC} = + r_{7} + r_{8} + r_{9} - 3r_{10} + r_{11} + r_{12} + r_{13}$$

$$\dot{c}_{A.BC} = + 1r_{10} - r_{11} - 2r_{14}$$

$$\dot{c}_{AB.C} = + 1r_{10} - r_{12} - 2r_{15}$$

$$\dot{c}_{ABC.} = + 1r_{10} - r_{13} - 2r_{16}$$
(9)

The three species of the trimer with one broken bond A.BC, AB.C and ABC. can be removed from the model using a steady state assumption:

– Lump the three states with one broken bond into the species XYZ. with balance equation in steady state

$$\dot{c}_{XYZ.} = +3k_2c_{ABC} - \bar{k}_1c_{XYZ.} - 2k_2c_{XYZ.} = 0.$$
⁽¹⁰⁾

- Now c_{XYZ} can be resolved as a function of c_{ABC} , namely:

$$c_{XYZ.} = \frac{3k_2}{\bar{k}_1 + 2k_2} c_{ABC}.$$
 (11)

- The decay of the state ABC is given by the new rate $r_{17} = -3r_{10} + \bar{k}_1 c_{XYZ}$. By inserting c_{XYZ} , which was calculated above, the decay rate becomes

$$r_{17} = -\left(3k_2 - \bar{k}_1 \frac{3k_2}{\bar{k}_1 + 2k_2}\right)c_{ABC} = \frac{2 \times 3k_2^2}{\bar{k}_1 + 2k_2}c_{ABC}$$
(12)

such that the effective decay rate constant of the trimer is

$$k_2^* = \frac{6k_2^2}{\bar{k}_1 + 2k_2}.\tag{13}$$

The solution of the resulting reduced ODE System

$$\dot{c}_{A} = -r_{1} + r_{2} - r_{5} + r_{6} - r_{7} + r_{17}/3$$

$$\dot{c}_{B} = -r_{1} + r_{2} - r_{3} + r_{4} - r_{8} + r_{17}/3$$

$$\dot{c}_{C} = +r_{3} - r_{4} - r_{5} + r_{6} - r_{9} + r_{17}/3$$

$$\dot{c}_{AB} = +r_{1} - r_{2} - r_{9} + r_{17}/3$$

$$\dot{c}_{BC} = +r_{3} - r_{4} - r_{7} + r_{17}/3$$

$$\dot{c}_{CA} = +r_{5} - r_{6} - r_{8} + r_{17}/3$$

$$\dot{c}_{ABC} = +r_{7} + r_{8} + r_{9} - r_{17}$$
(14)

by Matlab shows only marginal differences if compared with the full system (9).